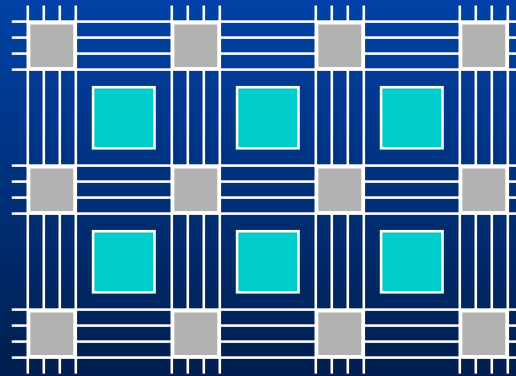


Processor and Integrated System Design in FPGAs

Jan Gray

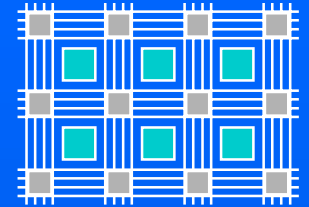
President, Gray Research LLC

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lw r12, 4 (r3)

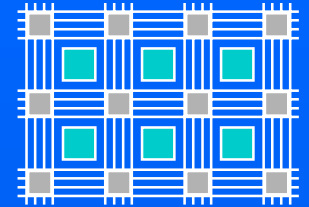
Introduction



lw r12,4(r3)

- ◆ Recent dense programmable logic enables practical “desktop processor development”
- ◆ Examples
 - ◆ RISC4005 (Freidin 1990)
 - ◆ j32 (1995)
 - ◆ XSOC/xr16 (1998)
 - ◆ xr16: simple pipelined 16-bit RISC, 33 MHz, 1.4 CPI
 - ◆ XSOC: on-chip bus and peripherals
- ◆ *Towards a reusable cores community and cost-effective FPGA-based systems-on-a-chip*

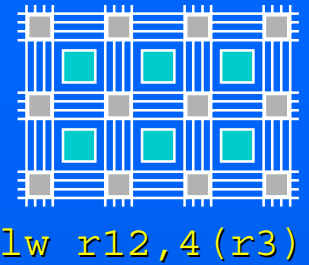
Outline



lw r12,4(r3)

- ◆ Introduction to FPGAs
 - ◆ Xilinx XC4000XL architecture
 - ◆ Development process
- ◆ XSOC/xr16 project
 - ◆ xr16 tools, architecture, datapath, control unit
 - ◆ XSOC design and implementation
- ◆ Ongoing work
- ◆ References, Resources

Field Programmable Gate Arrays

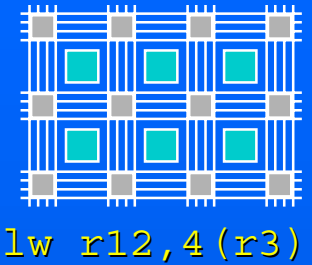


- ◆ Evolved from PALs, PLDs, Complex PLDs
- ◆ A niche hardware implementation technology

Technology	Gates	Speed	Cost	Part Cost	Spin time
Custom VLSI	<100M	<1 GHz	\$50K-up	\$1-up	weeks
Gate array	<10M	<400 MHz	\$10K-\$1M	\$1-up	days/weeks
FPGA	<1M	<200 MHz	\$100-\$100K	\$3-\$1K	minutes/hours

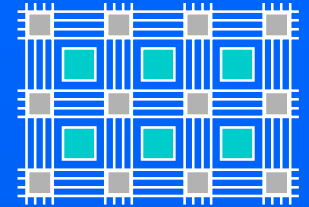
- ◆ Strengths: quick prototyping and time-to-market, reprogrammability, relatively easy to use
- ◆ Weaknesses: cost, density, speed
- ◆ Vendors: Xilinx, Altera, Actel, Atmel, Lucent, Cypress, Lattice, QuickLogic

FPGA Technologies



- ◆ Programmable Elements
 - ◆ Logic
 - ◆ combinatorial: lookup tables (LUTs) or gates + muxes
 - ◆ sequential: flip-flops
 - ◆ Interconnect: metal wire segments connected by...
 - ◆ pass transistors driven by SRAM or EEPROM bit cells
 - ◆ multiplexers
 - ◆ fuses or antifuses (one-time programmable)
- ◆ FPGA architecture variations
 - ◆ CLB arrays, rows, macrocells; fine/coarse grain
 - ◆ RAM, fast wide adders, TBUFs (3-state buffers)

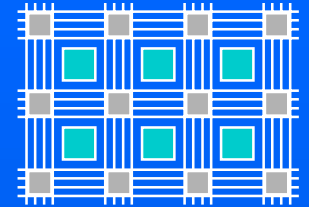
FPGAs: Easy to Use



lw r12,4(r3)

- ◆ Idealized digital design model
 - ◆ No analog EE: ignore gate and wire capacitance, transistor W/L ratios, etc.
 - ◆ Low skew global nets → No clock skew worries
 - ◆ Buffered line drivers → No gate fan-out worries
- ◆ Synchronous design “just works”
 - ◆ Avoid async. flip-flop state changes, gated clocks
 - ◆ Try to keep everything on one device
 - ◆ Design, simulate, it *should* just work
 - ◆ ...even if you’re just a software type
- ◆ But, effort to master effective use of resources

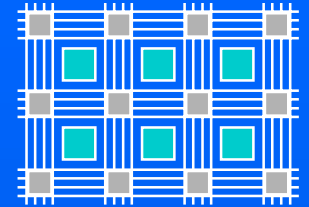
FPGA Applications



lw r12,4 (r3)

- ◆ Conventional: cheap, fast turnaround, field upgradable ASICs
 - ◆ Telecom, networking, XSOC/xr16
- ◆ (Re-) configurable computing
 - ◆ Quickturn: fast design simulation/verification
 - ◆ Signal, image processing: radar, radio
 - ◆ Graphics: compositing, octree rendering
 - ◆ Military: target dependent correlation/recognition
 - ◆ Cryptography: DES search
 - ◆ “Hardware” genetic algorithms

Xilinx XC4000 FPGAs

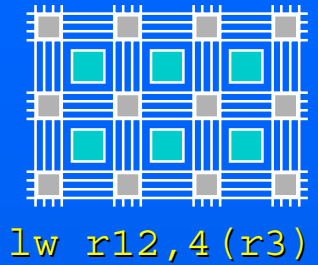


1w r12,4 (r3)

- ◆ Introduced in 1992, Xilinx's third generation
- ◆ SRAM based
- ◆ Device family spans 3,000 to 250,000 “gates”,
1 to 3 ns LUT delays, \$10 to \$1000 Q1
- ◆ XSOC uses 3.3V XC4005XLPC84-3
 - ◆ 14x14 Configurable Logic Blocks, 1.6 ns LUT delay
 - ◆ 60 bonded-out I/O Blocks
 - ◆ \$20 (\$3 Q100K)
- ◆ Eclipsed by newer Virtex, Virtex-E families
 - ◆ 256x16 SRAM blocks, up to 2M “gates”

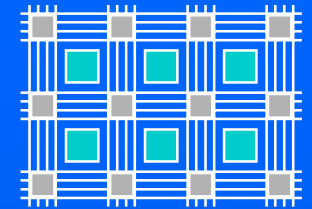
XC4000

Configurable Logic Block



- ◆ Two 4-input LUTs
- ◆ Two D flip-flops w/ clock enable
- ◆ Special features:
 - ◆ Writeable LUTs provide 32x1 or 16x2 SRAM
 - ◆ 8 CLBs make a 16x16 register file
 - ◆ Dedicated fast ripple carry chain hardware
 - ◆ 8 CLBs implement 16-bit adder, ~10 ns delay
- ◆ *Fig. 1*

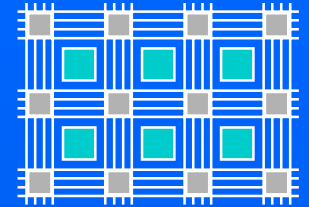
XC4000 I/O Block



lw r12,4(r3)

- ◆ Input: direct, registered
- ◆ Output: direct, registered, 3-stated, pull-ups/pull-downs, slew rate options
- ◆ *Fig. 11*

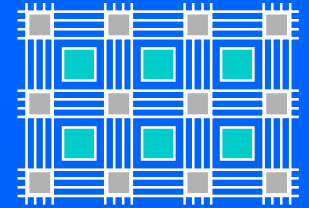
XC4000 Interconnect



lw r12,4 (r3)

- ◆ Programmable Interconnect Points
- ◆ Switch matrices
- ◆ Input multiplexers
- ◆ Hierarchical resources for hierarchical circuits
- ◆ Special features
 - ◆ Long lines - buses and control lines
 - ◆ 3-state buffer per output
 - ◆ Low skew global clock lines
- ◆ *Fig. 12-16*

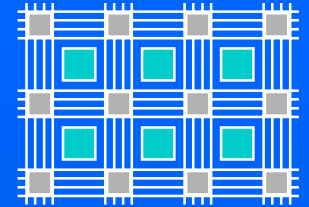
FPGA Development Process



`lw r12,4(r3)`

- ◆ Design capture: schematics, custom tools, hardware design languages; *floorplan!*
- ◆ Simulate/verify
- ◆ Compile: “technology mapping”, place, route
- ◆ Make config bitstream: XC4005XL: 152K bits
- ◆ Download:
 - ◆ “master load” from byte-wide or bit-serial ROM
 - ◆ “slave load” by microprocessor or XChecker pod
- ◆ Test/Debug; pod can readback internal state

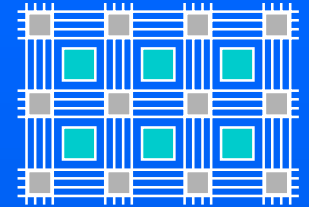
Talk Outline



lw r12,4(r3)

- ◆ Introduction to FPGAs
 - ◆ Xilinx XC4000XL architecture
 - ◆ Development process
- ◆ XSOC/xr16
 - ◆ xr16 tools, architecture, datapath, control unit
 - ◆ XSOC design and implementation
- ◆ Demo
- ◆ Ongoing work
- ◆ References, Resources

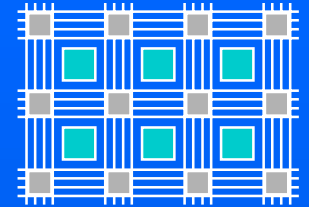
XSOC/xr16 Project Goals



1w r12,4 (r3)

- ◆ *Make integrated system dev more accessible*
- ◆ Subject of *Circuit Cellar* article series
 - ◆ Explain FPGA processor and system design
 - ◆ Build a fast, simple, thrifty, real CPU
 - ◆ Help launch a free cores development community
 - ◆ Provide on-chip bus architecture for easy cores reuse
 - ◆ Show FPGA CPUs can be practical (=cheap)
 - ◆ Not just an ASIC prototyping tool
 - ◆ Show C compiler support
 - ◆ Show efficient FPGA design “best practices”

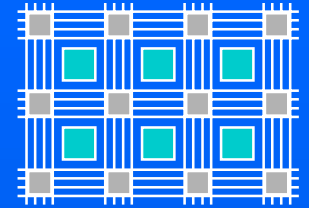
Development Tools



`lw r12, 4(r3)`

- ◆ Need C compiler for our nascent 16-bit RISC:
port LCC [Fraser and Hanson]
 - ◆ `mips.md` → `xr16.md`
 - ◆ 32 registers → 16 registers
 - ◆ `sizeof(int) == sizeof(void*) == 4` → 2
 - ◆ Misc: branches; long ints; software mul/div
- ◆ Assembler
 - ◆ Lex, parse, fix far branches, apply fixups, emit
- ◆ Instruction set simulator

LCC Generated Code



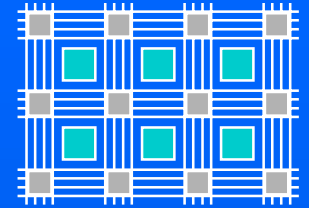
lw r12,4(r3)

```
◆ typedef struct TN {  
    int k;  
    struct TN *left, *right;  
} *T;
```

```
T search(int k, T p) {  
    while (p && p->k != k)  
        if (p->k < k)  
            p = p->right;  
        else  
            p = p->left;  
    return p;  
}
```

```
◆ _search:                ; r3=k r4=p  
    br L3  
L2: lw r9,(r4)  
    cmp r9,r3              ; p->k < k?  
    bge L5  
    lw r4,4(r4)           ; p=p->right  
    br L6  
L5: lw r4,2(r4)           ; p=p->left  
L6:  
L3: mov r9,r4  
    cmp r9,r0              ; p==0?  
    beq L7  
    lw r9,(r4)  
    cmp r9,r3              ; p->k != k?  
    bne L2  
L7: mov r2,r4              ; retval=p  
L1: ret
```

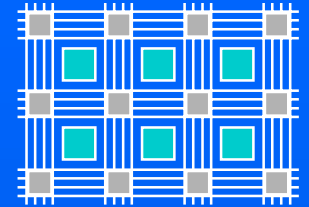

XR16 Instruction Set Design



`lw r12, 4(r3)`

- ◆ Goals and constraints
 - ◆ Compact 16-bit instructions
 - ◆ 16 16-bit registers
 - ◆ Cover integer C operations
 - ◆ KISS
 - ◆ Fast → pipelined → regular
 - ◆ Byte addressable → load/store bytes/words
 - ◆ `disp(reg)` addressing
 - ◆ Long ints → add with carry, shift extended
 - ◆ Scale to 32-bit registers
- ◆ Common instructions
 - ◆ `lw sw mov lea call br cmp`
 - ◆ `mov lea cmp` → add/sub
 - ◆ `disp(reg)` addressing (69%)
- ◆ Resolved:
 - ◆ 3 operand: add sub addi
 - ◆ 4-bit immediates, 12-bit immediate prefix
 - ◆ no condition codes; interlocked add/sub+*bcond*
 - ◆ fast call/return → call/jal
 - ◆ mul/div/shifts in software

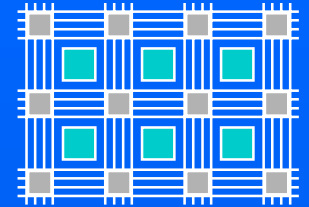
XR16 Instruction Set



`lw r12,4(r3)`

Hex	Fmt	Assembler	Semantics	N
0dab	rrr	add rd,ra,rb	rd = ra + rb;	1
1dab	rrr	sub rd,ra,rb	rd = ra - rb;	1
2dai	rri	addi rd,ra,imm	rd = ra + imm;	1
3d*b	rr	{and or xor andn adc sbc} rd,rb	rd = rd <i>op</i> rb;	1
4d*i	ri	{andi ori xori andni adci sbci slli slxi srai srli srxi} rd,imm	rd = rd <i>op</i> imm;	1
5dai	rri	lw rd,imm(ra)	rd = *(int*)(ra+imm);	2
6dai	rri	lb rd,imm(ra)	rd = *(byte*)(ra+imm);	2
8dai	rri	sw rd,imm(ra)	*(int*)(ra+imm) = rd;	2
9dai	rri	sb rd,imm(ra)	*(byte*)(ra+imm) = rd;	2
Adai	rri	jal rd,imm(ra)	rd = pc, pc = ra + imm;	3
B*dd	br	{br brn beq bne bc bnc bv bnv blt bge ble bgt bltu bgeu bleu bgtu} label	if (<i>cond</i>) pc += 2*disp8;	*
Ciii	i12	call func	r15 = pc, pc = imm12<<4;	3
Diii	i12	imm imm12	imm'next _{15:4} = imm12;	1
7xxx Exxx Fxxx	-	<i>reserved</i>	<i>reserved</i>	

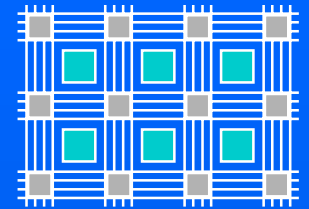
Synthesized Instructions



lw r12,4(r3)

Instruction	Maps to
nop	and r0,r0
mov rd,ra	add rd,ra,r0
cmp ra,rb	sub r0,ra,rb
subi rd,ra,im	addi rd,ra,-im
cmpi ra,im	addi r0,ra,-im
com rd	xori rd,-1
lea rd,imm(ra)	addi rd,ra,imm
lbs rd,imm(ra) <i>(load-byte, sign-extending)</i>	lb rd,imm(ra) xori rd,0x80 subi rd,0x80
j addr	jal r0,addr
ret	jal r0,0(r15)

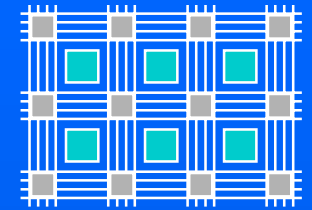
Compromised Instructions



`lw r12, 4(r3)`

- ◆ No multiply/divide
- ◆ No barrel shifter: 1-bit shifts only
- ◆ Jumps and taken branches take 3 cycles
 - ◆ Annul the two branch shadow instructions
- ◆ Loads/stores take at least 2 cycles
 - ◆ One shared memory port
- ◆ Simple interrupt/return
- ◆ No floating point

Design Hierarchy: xr16



lw r12,4(r3)

- ◆ XSOC
 - ◆ xr16
 - ◆ datapath
 - ◆ control unit
 - ◆ memctrl
 - ◆ vga
 - ◆ parin, parout, iram

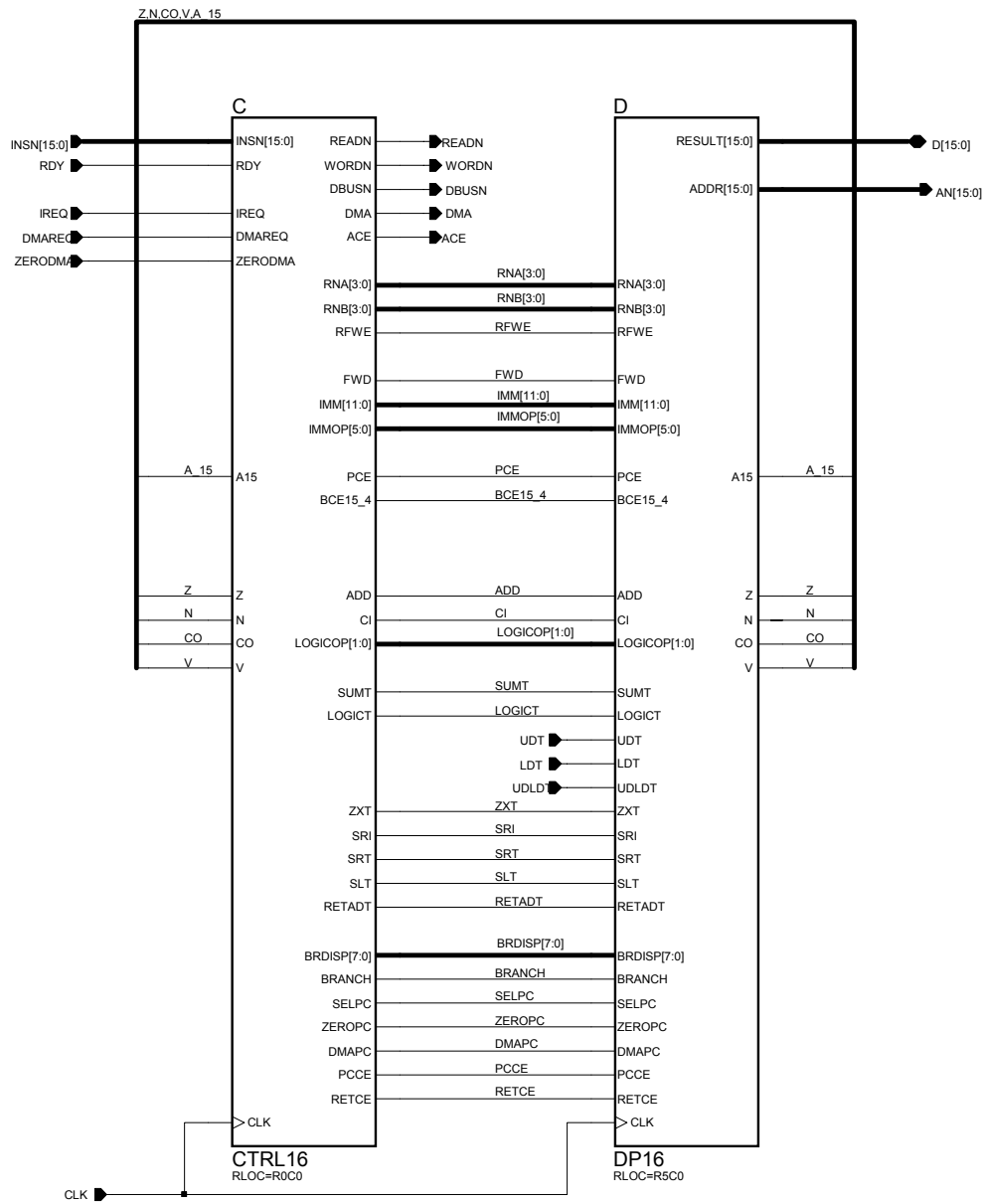
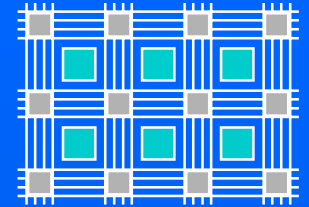


Figure S2: XR16 CPU Schematic

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This work and its use subject to XSOC License Agreement. See LICENSE file.	Macro: XR16
	Date: 02/23/100

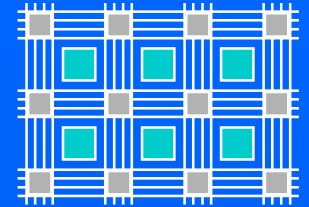
Datapath Resources



`lw r12, 4(r3)`

- ◆ 2 read, 1 write per cycle register file
- ◆ Immediate operand multiplexer
- ◆ A and B operand registers
- ◆ Adder, logic unit, shifter
- ◆ Result multiplexer
- ◆ Zero, negative, carry, overflow detection
- ◆ PC, branch displacement multiplexer, PC adder
- ◆ Address multiplexer

Design Hierarchy



lw r12,4(r3)

- ◆ XSOC
 - ◆ xr16
 - ◆ datapath
 - ◆ control unit
 - ◆ memctrl
 - ◆ vga
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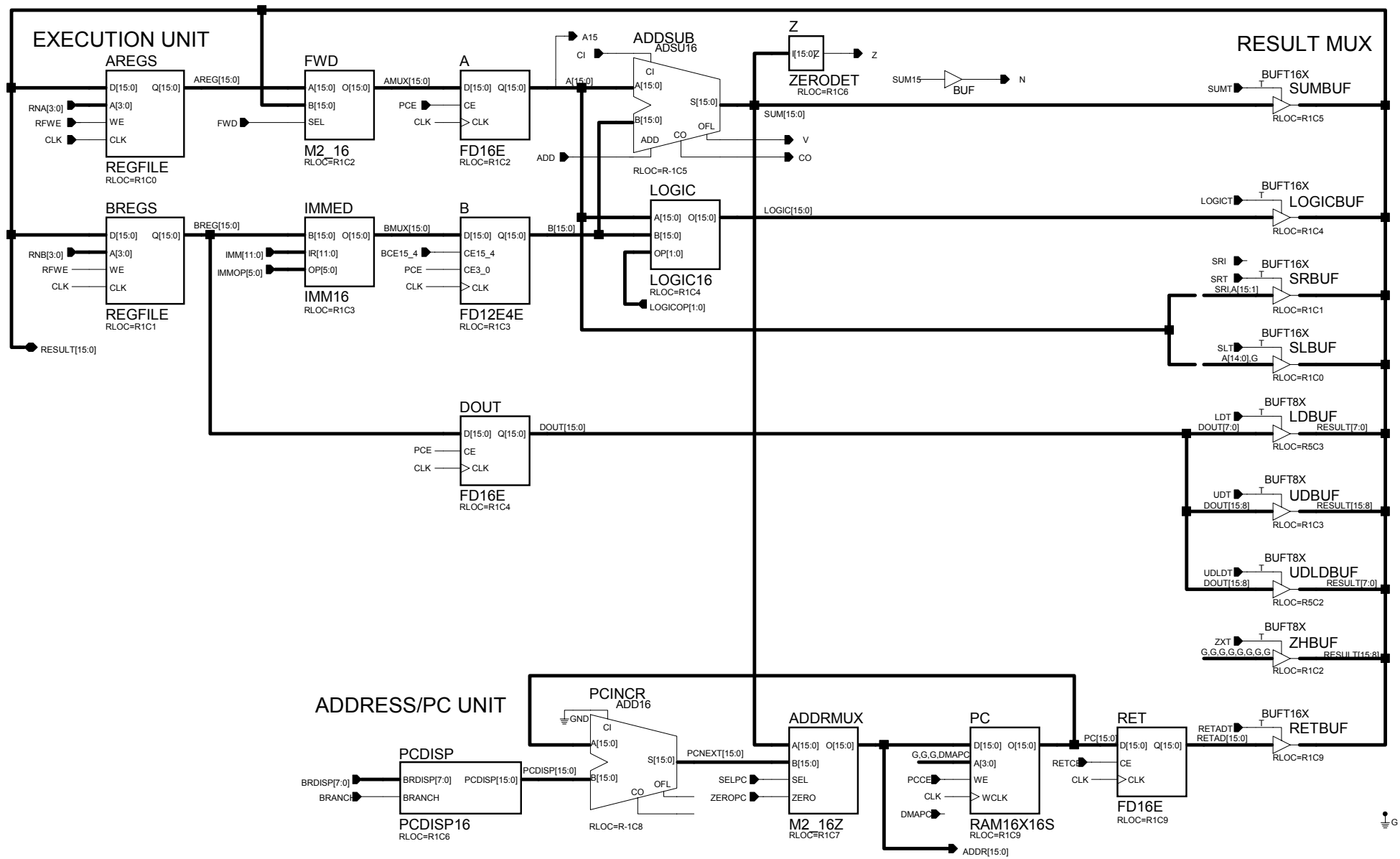
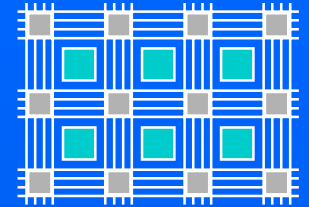


Figure S3: XR16 CPU Datapath Schematic

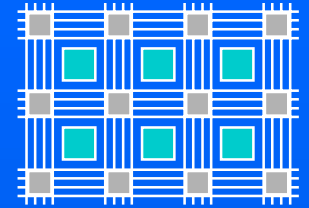
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	Date: 02/23/100

Control Unit: Pipeline Stages



lw r12, 4(r3)

- ◆ IF: Instruction Fetch
 - ◆ Read instruction at addr; prepare $\text{addr} \leftarrow \text{PC} += 2$; capture instruction in Instruction Register (IR)
- ◆ DC: Decode and register file access
 - ◆ Decode instruction; read register operands; prepare immediate field; select operands
- ◆ EX: Execute
 - ◆ Add, logic, shift; select result; write to register file
 - ◆ call/jal: $\text{PC} := \text{sum}$
 - ◆ load/store: $\text{addr} \leftarrow \text{sum}$; stop pipe for access

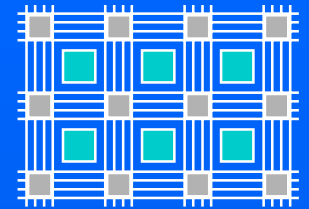


`lw r12, 4(r3)`

Control Unit: Pipeline Hazards

- ◆ Memory wait states
- ◆ Result use data dependence
- ◆ Load/store memory access
- ◆ Jumps and taken branches
- ◆ Interrupts

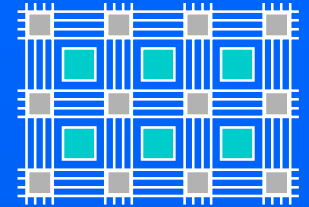
Control Unit Elements



`lw r12,4(r3)`

- ◆ Control finite state machine
 - ◆ Memory interface controls
- ◆ Instruction registers (next, DC stage, EX stage)
- ◆ DC: Instruction decoder
- ◆ DC: Register file read controls
- ◆ DC: Immediate operand and forwarding control
- ◆ EX: ALU and result multiplexer control
- ◆ EX: Register file write control
- ◆ DC+EX: Conditional branch control

Design Hierarchy



lw r12,4(r3)

- ◆ XSOC
 - ◆ xr16
 - ◆ datapath
 - ◆ control unit
 - ◆ memctrl
 - ◆ vga
 - ◆ parin, parout, iram

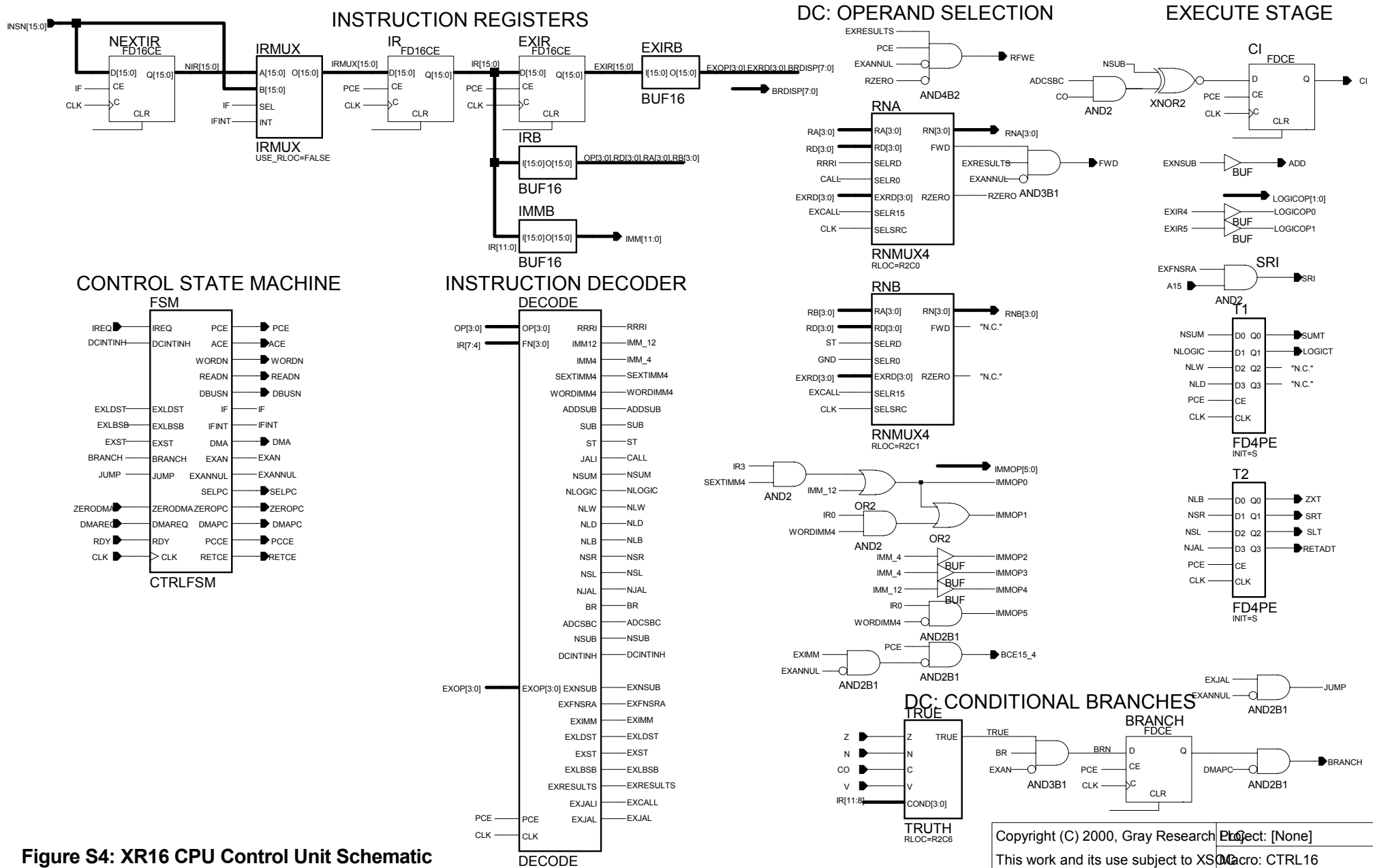
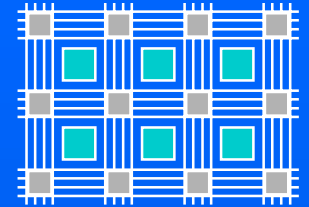


Figure S4: XR16 CPU Control Unit Schematic

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 Date: 02/23/100

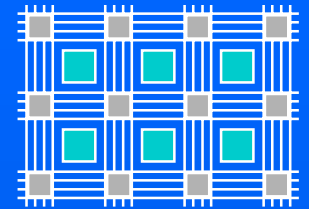
Outline



lw r12,4(r3)

- ◆ Introduction to FPGAs
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 - ◆ XSOC design and implementation
- ◆ Ongoing work
- ◆ References, Resources

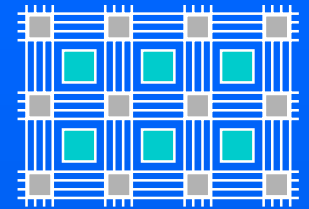
XSOC Project



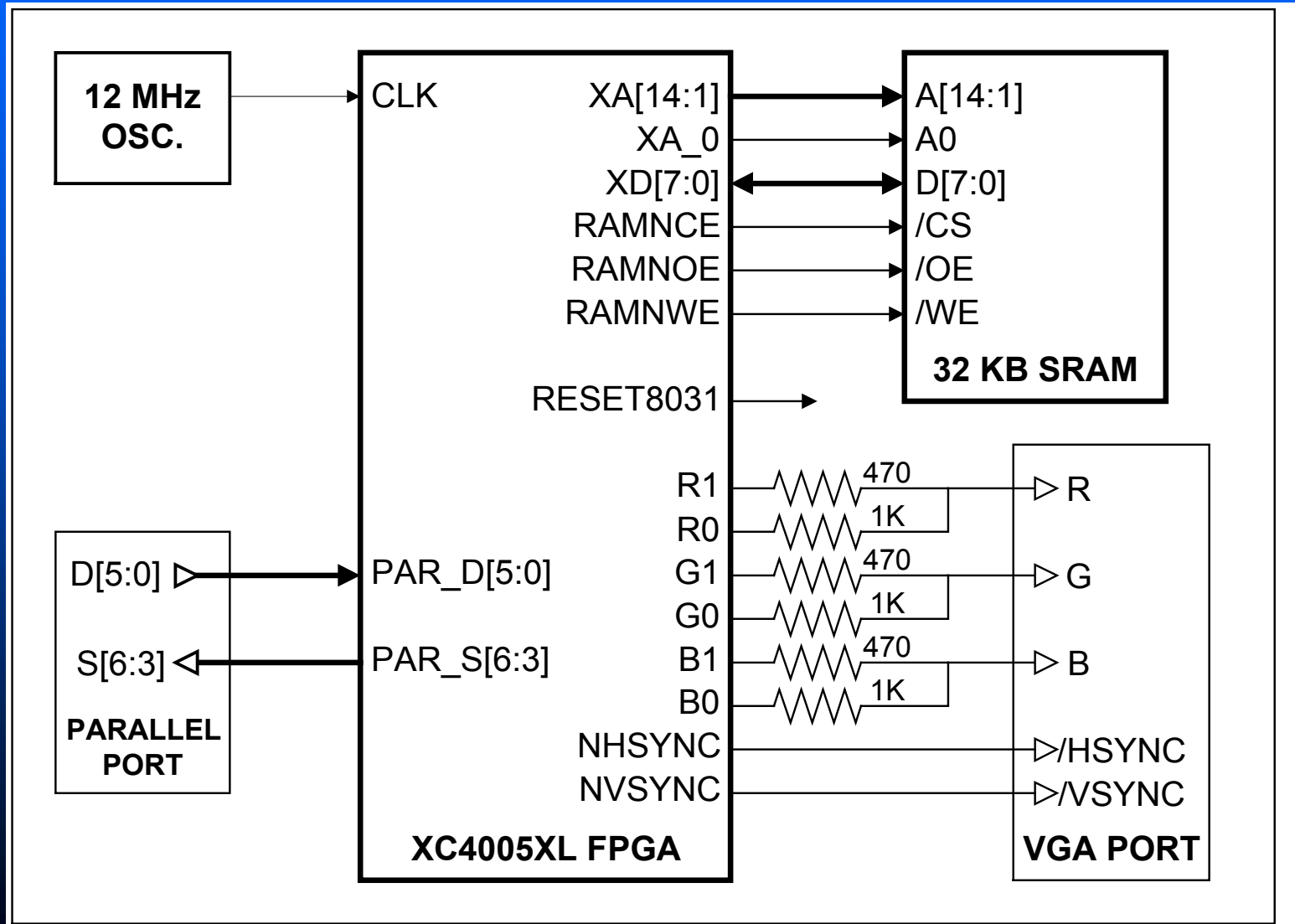
lw r12,4(r3)

- ◆ *Make integrated system dev more accessible*
- ◆ Build with Xilinx Student Edition (\$100)
- ◆ Target XESS XS40-005XL student lab board
 - ◆ Textbook, exercises, and tools (\$100)
- ◆ Integrated system:
 - ◆ Processor, memory, memory controller, on-chip bus, parallel port, ram, video controller
 - ◆ Double-cycle byte-wide external SRAM
 - ◆ 32 KB SRAM → 576x455 monochrome display

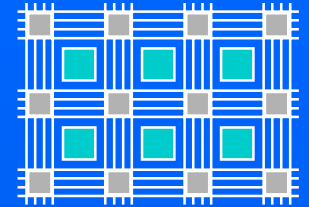
“System-on-a-Chip” in Context



r12,4 (r3)

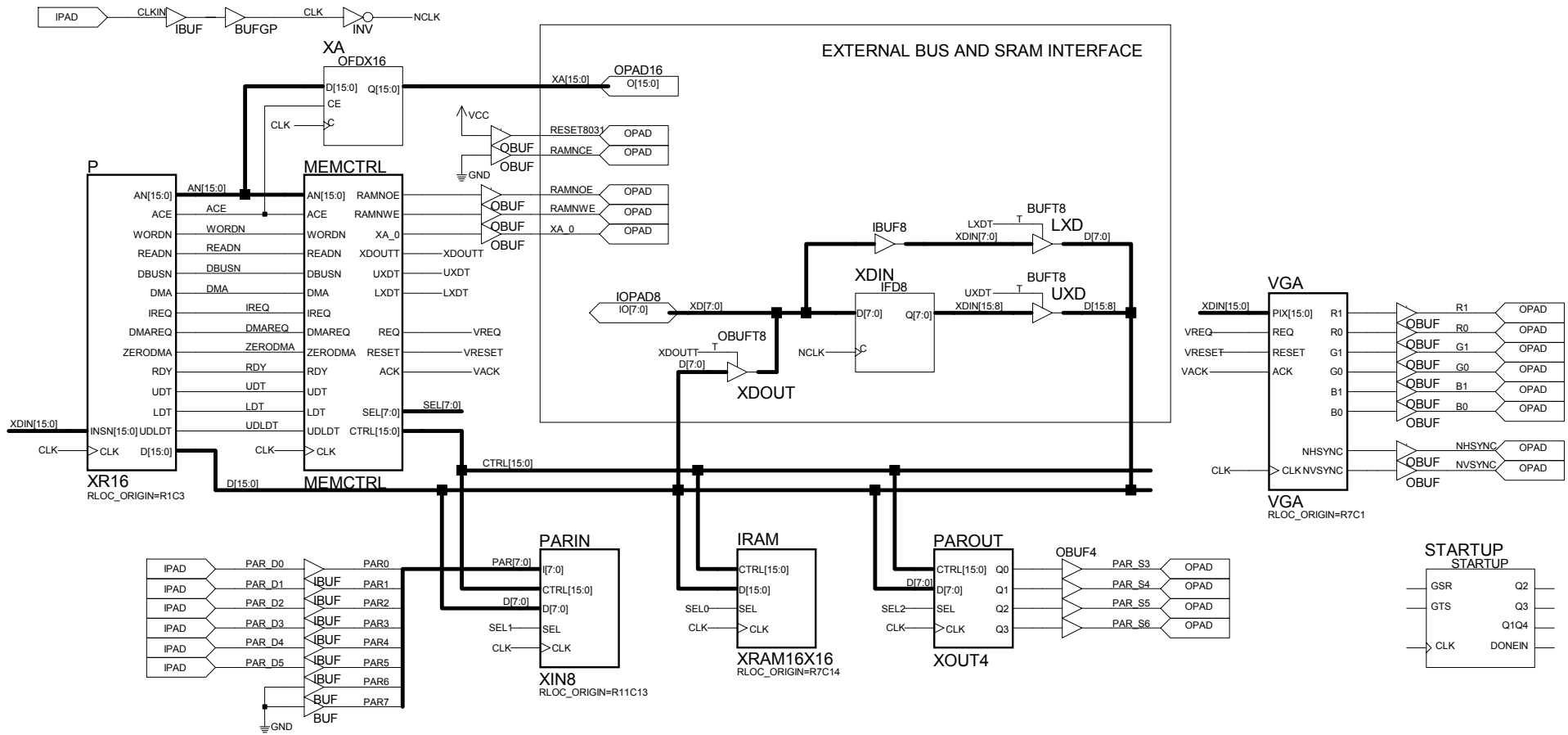


Design Hierarchy

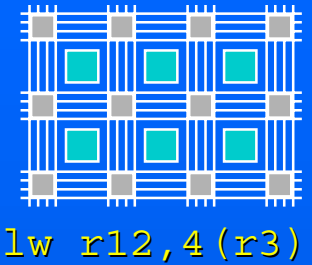


`lw r12,4(r3)`

- ◆ XSOC
 - ◆ xr16
 - ◆ datapath
 - ◆ control unit
 - ◆ memctrl
 - ◆ vga
 - ◆ parin, parout, iram

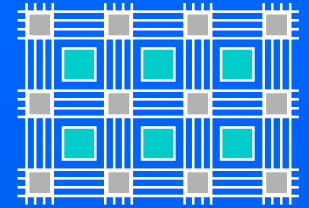


On-Chip Bus



- ◆ 16-bit pipelined on-chip data bus
- ◆ Bus/memory controller
 - ◆ Address decoding, bus controls (output enables, clock enables), external SRAM control
- ◆ Make core reuse easy: *no glue logic required*
 - ◆ Abstract control signal bus
 - ◆ Locally decoded within each core
 - ◆ Just add core, attach data, control, and select lines
 - ◆ Can evolve bus with new features without invalidating existing designs and cores

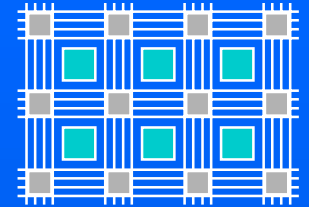
Video Controller



lw r12,4(r3)

- ◆ 32 KB → 576x455 monochrome bitmap
- ◆ VGA compatible sync timings
- ◆ Video signal generation
 - ◆ A series of frames, each a series of lines, each a series of pixels
 - ◆ Fetch 16-pixel words, shifts them out serially
 - ◆ DMA read a new word every 8 clocks
 - ◆ Drive horizontal and vertical sync to “frame” pixels
 - ◆ 2 10-bit counters and 4 10-bit comparators

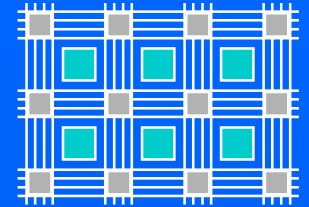
System Bring Up



`lw r12,4(r3)`

- ◆ 11/98: Proposal, compiler, instruction set, datapath
- ◆ 12/98: Control unit, simulate CPU, target XS40 board, 1 Hz, 20 MHz
- ◆ Debugging with LEDs and stop: goto stop;
- ◆ Later added on-chip bus, external SRAM interface, DMA, video, and interrupts
- ◆ Testing challenge

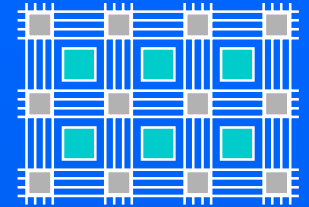
Recap



`lw r12,4(r3)`

- ◆ Final floorplan
- ◆ FPGA editor
- ◆ Demo
 - ◆ Source
 - ◆ Listing
 - ◆ Compile and go

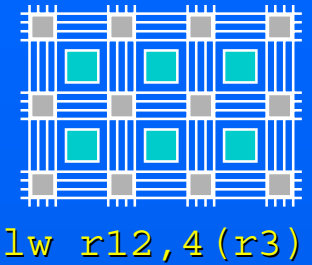
Ongoing Work



lw r12,4(r3)

- ◆ Near term
 - ◆ Package XSOC/xr16 distribution for beta test ✓
 - ◆ Retarget to Verilog ✓
 - ◆ Port to Virtex and Altera architectures
- ◆ Longer term
 - ◆ Help build a community
 - ◆ xr32 $\frac{1}{2}$ ✓
 - ◆ Port GCC/binutils; build eCOS, Linux?
 - ◆ Chip multiprocessors

References



◆ FPGAs

- ◆ Trimberger, S., *Field-Programmable Gate Array Technology*, Kluwer.

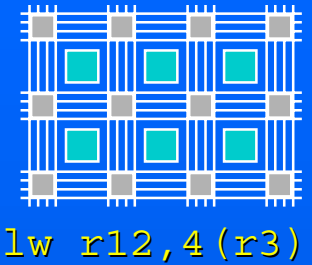
◆ LCC

- ◆ C. Fraser and D. Hanson, *A Retargetable C Compiler: Design and Implementation*, Benjamin/Cummings.

◆ RISC architecture/implementation

- ◆ D. Patterson and J. Hennessy, *Computer Organization and Design: The Hardware/Software Interface*, also *Computer Architecture: A Quantitative Approach*, Morgan Kaufmann

Resources



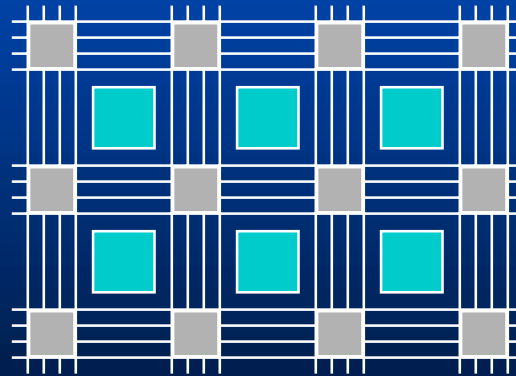
- ◆ Usenet: comp.arch.fpga, comp.arch
- ◆ Web
 - ◆ www.xilinx.com/products/xc4000XLA.htm
 - ◆ www.io.com/~guccione/HW_list.html
 - ◆ www.fpgacpu.org
- ◆ Conferences
 - ◆ SIGDA Int'l Symp. on FPGAs, Monterey, Feb.
 - ◆ IEEE Symp. on FPGAs for Custom Computing Machines, Napa, April

Processor and Integrated System Design in FPGAs

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`lw r12, 4(r3)`